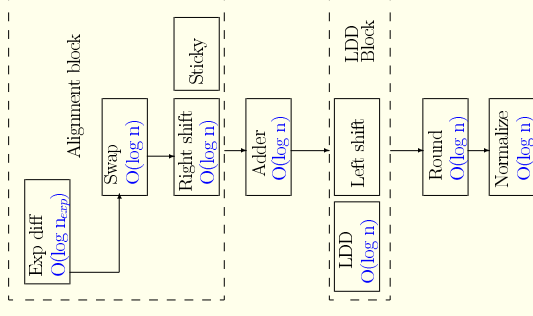


Computer Arithmetic, Lecture 12:
Recent adders

Hossam A. H. Fahmy

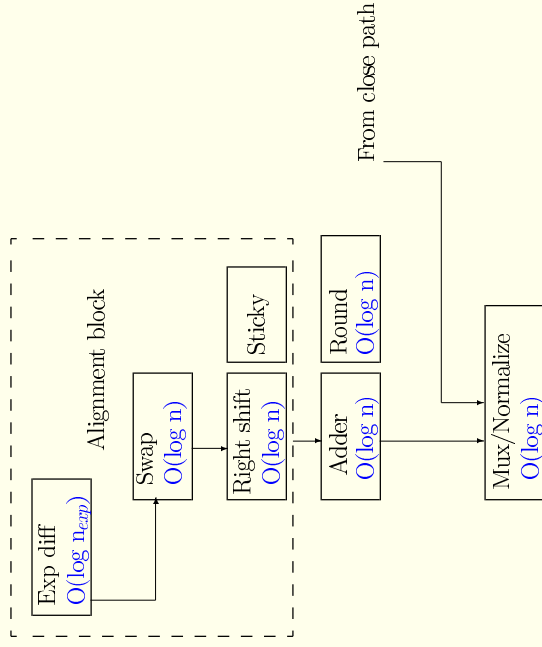


Once again the one-path



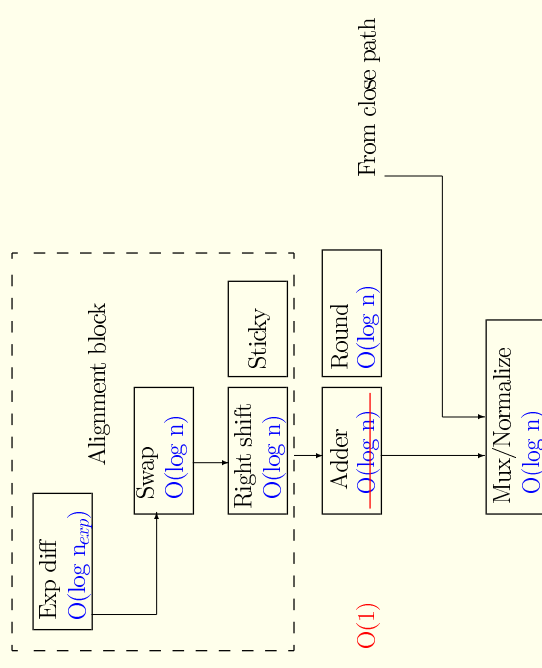
Time delays in the blocks of an adder (one-path algorithm)

The improved two-path



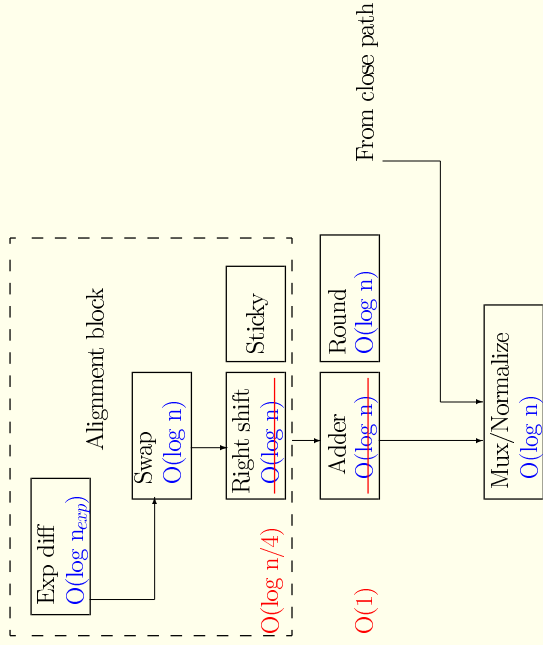
Time delays in the blocks of an adder (two-path algorithm)

Now comes redundancy



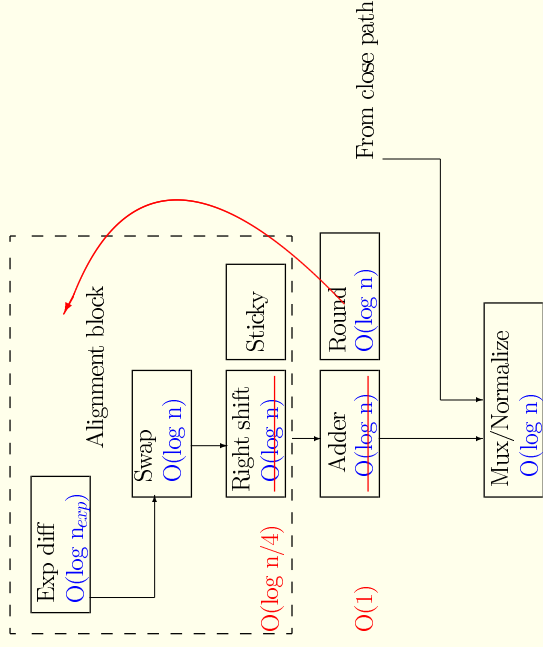
Time delays with redundancy

Now comes redundancy



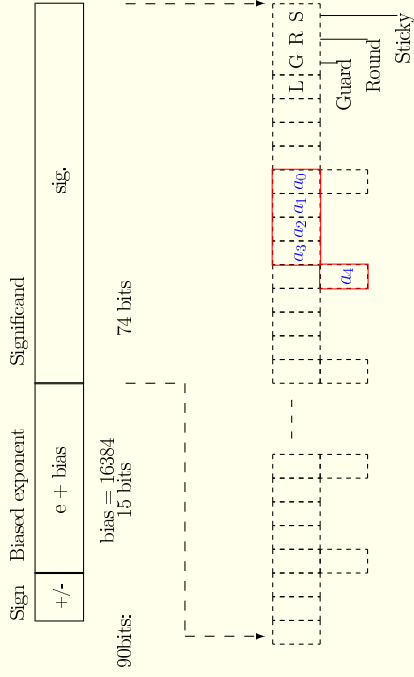
Time delays with redundancy

Now comes redundancy



Time delays with redundancy

The representation

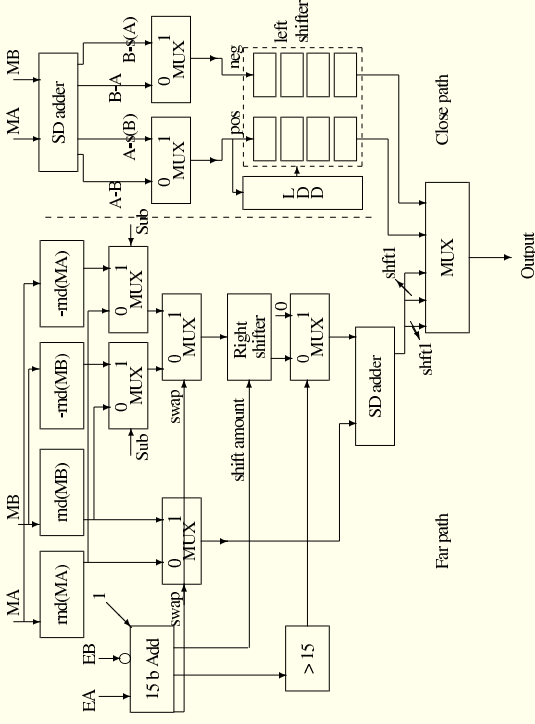


The proposed SD format for floating point numbers

The gain

- Normalized numbers in the IEEE format are (almost) a subset of the presented format.
- The conversion back to the IEEE format takes place only when a register is to be stored in the main memory and the conversion delay is overlapped with the store operation.

Therefore, contrary to previous designs using SD numbers, the proposed system effectively hides the conversion delay.



General blocks of the proposed adder (two-path algorithm)

Circuit statistics and simulation results

	FP-Adder	FP-Multiplier
nodes	46845	76523
NMOS	63589	104695
PMOS	61649	105037
Test vectors	5000	10000
Model delay	34FO4	35FO4
Sim delay(0.6μm)	14ns	14.8ns
Sim delay(0.3μm)	(33.35FO4)	(35.25FO4)
	6ns	6.4ns
	(32.40FO4)	(34.60FO4)

Peculiar Leading digits

For regular binary, if the result starts by few zeros like 00110... we need to shift that result to the left in order to normalize it.

Beside that simple case, in the current design we might get:

0 1 -15 ... -15 l ... = 0 0 0 0 ... 1 l ...
 0 -1 15 ... 15 l ... = 0 0 0 0 ... -1 l ...

An extension of N and P recoders, originally proposed by Daumas and Matula in 1997, is used to correctly find the leading digit in such cases.

Cost?

	n = 60	n = 80	n = 120	
4 bits per digit	width increase(%)	33.3	31.6	29.6
	speed up (%)	10.5	12.5	16.7
8 bits per digit	width increase(%)	29.3	21.1	18.5
	speed up (%)	7.9	10	14.3

Floating point adder trade off when the fan-in limit is 3.

Another FP redundant adder

- The addition is the most frequent operation.
- Two subsequent *dependent* additions will take a long time on a pipelined machine.
- Nielsen *et al.* forward a redundant result to the following addition after just two clock cycles.
- In the third clock cycle, the rounding decision is calculated and is forwarded to the dependent addition.
- The fourth clock cycle produces the full non-redundant result.

The packet forwarding adder

- The adder accepts one of the operands in a redundant form and in two “packets”.
- It is possible to start a new dependent addition every two clock cycles.
- The result is available, correctly rounded, in the IEEE format in four clock cycles.

A detailed design

Smith *et al.* in 1999 provided a good amount of details for their various designs.

Their paper is instructive to teach you about the trade-off in considering the amounts of shifts needed and the decision on which is the largest significant.

This is one of a few papers that speak about the exponents, the exceptional inputs, and shows a real layout for a fabricated chip.

A comparative analysis

Seidel and Even in 2001 presented a nice comparative analysis of a large number of adders from the literature.

Although it is not exhaustive (and has minor mistakes such as in reference to work of Smith mentioned earlier) it is a very good source to understand the various trade-offs.

Real world issues

- The subnormal numbers are not easy to handle. For those interested, see the corresponding paper.
- Some machines must support multiple formats. The IBM mainframes still provide the old hexadecimal side by side with the binary.
- The need to decimal FP is growing and that too should be supported.

To sum up the addition

- Addition is very important and hence it received the attention of too many people over time.
- There is still a possibility to improve, especially when it comes to energy consumption.
- Another direction for future research is the support of multiple formats: binary, hexadecimal, and decimal.